# Root Finding and Optimization 

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## Root Finding

- It is not uncommon in scientific computing to want solve an equation numerically.
- If there is one unknown and one equation only, we can always write the equation as

$$
f(x)=0
$$

- If $x$ satisfies this equation, it is called a "root".
- If there's a set of equations, one can write:

$$
\begin{aligned}
& \mathbf{f}_{1}\left(x_{1}, x_{2}, x_{3}, \ldots\right)=0 \\
& f_{2}\left(x_{1}, x_{2}, x_{3}, \ldots\right)=0 \\
& \mathbf{f}_{3}\left(x_{1}, x_{2}, x_{3}, \ldots\right)=0
\end{aligned}
$$

The one-dimensional case is considerably easier to solve: First.

## 1D root finding



Algorithms always start from an initial guess and (usually) a bounding interval [a, b] in which the root is to be found.

What's so nice about 1D?

- If $\mathbf{f}(\mathbf{a})$ and $\mathbf{f}(\mathbf{b})$ have opposite signs, and $\mathbf{f}$ is continuous, there must be a root inside the interval: the root is bracketed.
- Consecutive refinement of the interval until $\mathbf{a}-\mathbf{b}<\varepsilon$ guarranteed to find the root.


## Bracketing

Must find bounding interval first.
Strategies:

- Plot the function.
- Make a conservative guess for [a, b] then slice up the interval, checking for sign change of $\mathbf{f}$.
- Make a wild guess and expand the interval until f exhibits a sign change.

Troublemakers


No sign change


Easily missed


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## Suppose we've bracketed a root, what's next?

Classic root finding algorithms

- Bisection
- Secant/False Position
- Ridders'/Brent
- Newton-Raphson (requires derivatives)

All of these zero in on the root, but have different convergence and stability characteristics.

Note:

- For polynomial f: specialized routines (Muller, Laguerre)
- For eigenvalue problems: specialized routines (linear algebra)
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## Bisection method

Split the interval. Compute $\mathbf{f}$ (midpoint). Get new root bracket.


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## False Position Method

Linearly interpolate $\mathbf{f}$. Solve for next $\mathbf{x}$. Get new root bracket.


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## Other non-derivate methods

- Secant: Like false position, but keep most-recent point.
- Ridders' method: uses an exponential fit.
- Brent's method: uses an inverse quadratic fit.


## Newton-Raphson method

Guess $\mathbf{x}$. From function and derivate, approximate root. Repeat.


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Guess $\mathbf{x}$. From function and derivate, approximate root. Repeat.


## Convergence and Stability

| method | convergence | stability |
| :--- | :--- | :--- |
| Bisection | $\epsilon_{\mathbf{n}+1}=\frac{1}{2} \epsilon_{\mathbf{n}}$ | Stable |
| Secant | $\epsilon_{\mathbf{n}+1}=\mathbf{c} \epsilon_{\mathbf{n}}^{1.6}$ | No bracket guarrantee |
| False position | $\epsilon_{\mathbf{n}+\mathbf{1}}=\frac{1}{2} \epsilon_{\mathbf{n}}-\mathbf{c} \epsilon_{\mathbf{n}}^{1.6}$ | Stable |
| Ridders' | $\boldsymbol{\epsilon}_{\mathbf{n}+2}=\mathbf{c} \epsilon_{\mathbf{n}}^{2}$ | Stable |
| Brent | $\boldsymbol{\epsilon}_{\mathbf{n}+\mathbf{1}}=\frac{1}{2} \epsilon_{\mathbf{n}}-\mathbf{c} \epsilon_{\mathbf{n}}^{\mathbf{2}}$ | Stable |
| Newton-Raphson | $\boldsymbol{\epsilon}_{\mathbf{n}+\mathbf{1}}=\mathbf{c} \epsilon_{\mathbf{n}}^{2}$ | Can be unstable |

## GNU Scientific Library (GSL)

Is a C library containing many useful scientific routines:

- Root finding
- Minimization
- Sorting
- Linear algebra, Eigen-systems
- Fast Fourier transforms
- Integration, differentiation, interpolation, approximation
- Random numbers
- Statistics, histograms, fitting
- Monte Carlo integration, simulated annealing
- ODEs
- Polynomials, permutations
- Special functions
- Vectors, matrices


## GSL root finding example

```
#include <iostream>
#include <gsl/gsl_roots.h>
struct Params {
    double v, w, a, b, c;
};
double examplefunction(double x, void * param) {
    Params* p = (Params*)param;
    return p->a*cos(sin(p->v+p->w*x))+p->b*x-p->c*x*x;
}
int main() {
    double x_lo = -4.0;
    double x_hi = 5.0;
    Params args = {0.3, 2/3.0, 2.0, 1/1.3, 1/30.0};
    gsl_root_fsolver* solver;
    gsl_function fwrapper;
    solver=gsl_root_fsolver_alloc(gsl_root_fsolver_brent);
    fwrapper.function = examplefunction;
    fwrapper.params = &args;
    gsl_root_fsolver_set(solver, &fwrapper, x_lo, x_hi);
```


## GSL root finding example

```
    std::cout << " iter [ lower, upper] root err\n";
    int status = 1;
    for (int iter=0; status and iter < 100; ++iter) {
    gsl_root_fsolver_iterate(solver);
    double x_rt = gsl_root_fsolver_root(solver);
    double x_lo = gsl_root_fsolver_x_lower(solver);
    double x_hi = gsl_root_fsolver_x_upper(solver);
    std::cout << iter <<' '<< x_lo <<' '<< x_hi
    <<' '<< x_rt <<' '<< x_hi - x_lo << "\n";
    status = gsl_root_test_interval(x_lo,x_hi,0,0.001);
}
gsl_root_fsolver_free(solver);
return status;
```

\}

## GSL root finding example

```
    std::cout << " iter [ lower, upper] root err\n";
    int status = 1;
    for (int iter=0; status and iter < 100; ++iter) {
        gsl_root_fsolver_iterate(solver);
        double x_rt = gsl_root_fsolver_root(solver);
        double x_lo = gsl_root_fsolver_x_lower(solver);
        double x_hi = gsl_root_fsolver_x_upper(solver);
    std::cout << iter <<' '<< x_lo <<' '<< x_hi
    <<' '<< x_rt <<' '<< x_hi - x_lo << "\n";
    status = gsl_root_test_interval(x_lo,x_hi,0,0.001);
}
gsl_root_fsolver_free(solver);
return status;
```

\}

Compilation and likage:

```
$ g++ -c -02 -I GSLINCDIR gslrx.cc -o gslrx.o
$ g++ gslrx.o -o gslrx -L GSLLIBDIR -lgsl -lgslcblas
```

(specify directories if installed in non-standard locations)

## Multidimensional Root Finding

$$
\overrightarrow{\mathbf{f}}(\overrightarrow{\mathrm{x}})=\overrightarrow{\mathbf{0}} \quad \text { or } \quad \begin{aligned}
\mathrm{f}_{1}\left(\mathrm{x}_{1}, x_{2}, x_{3}, \ldots, x_{D}\right) & =0 \\
\mathrm{f}_{2}\left(\mathrm{x}_{1}, x_{2}, x_{3}, \ldots, x_{D}\right) & =0 \\
& \vdots \\
& \\
\mathrm{f}_{\mathrm{D}}\left(\mathrm{x}_{1}, x_{2}, x_{3}, \ldots, x_{D}\right) & =0
\end{aligned}
$$

- Cannot bracket a root with a finite number of points.
- Roots of each equation define a D - $\mathbf{1}$ hypersurface.
- Looking for possible intersections of hypersurfaces.


## Newton-Raphson for Multidimensional Root Finding

Given a good initial guess, Newton-Raphson can work in arbitrary dimensions:

$$
\begin{aligned}
\overrightarrow{\mathrm{f}}\left(\overrightarrow{\mathrm{x}}_{0}\right. & +\delta \overrightarrow{\mathrm{x}})=\overrightarrow{0} \\
\overrightarrow{\mathrm{f}}\left(\overrightarrow{\mathrm{x}}_{0}\right) & +\frac{\partial \overrightarrow{\mathrm{f}}}{\partial \overrightarrow{\mathrm{x}}} \cdot \delta \overrightarrow{\mathrm{x}}=\overrightarrow{0} \\
\overrightarrow{\mathrm{f}}\left(\overrightarrow{\mathrm{x}}_{0}\right) & =-\frac{\partial \overrightarrow{\mathrm{f}}}{\partial \overrightarrow{\mathrm{x}}} \cdot \delta \overrightarrow{\mathrm{x}} \\
\overrightarrow{\mathrm{f}}\left(\overrightarrow{\mathrm{x}}_{0}\right) & =-\mathrm{J} \cdot \delta \overrightarrow{\mathrm{x}} \\
\delta \overrightarrow{\mathrm{x}} & =-\mathrm{J}^{-1} \cdot \overrightarrow{\mathrm{f}}\left(\overrightarrow{\mathrm{x}}_{0}\right)
\end{aligned}
$$

Requires inverting a $\mathbf{D} \times \mathbf{D}$ matrix, or at least, solving a linear set of equations: see lecture on linear algebra.
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## Convergence and Stability

- As in 1D, Newton-Raphson can be unstable.
- Need some safe guard that our iteration steps do not spin out of control.
- Several ways, e.g. make sure that $\|\overrightarrow{\mathbf{f}}(\overrightarrow{\mathbf{x}})\|^{2}$ gets smaller in each time step.
- This can potentially still fail, but usually does the trick.


## Optimization

## Optimization

Optimization $=$ Minimization $=$ Maximization
There's a function $\mathbf{f}\left(\mathbf{x}_{1}, \mathbf{x}_{2}, \ldots, \mathbf{x}_{\mathrm{D}}\right)$.

- Want to know for which set of $\mathbf{x}_{1}, \mathrm{x}_{2}, \ldots$ the function is at its minimum.
- Maximizing $\mathbf{f}$ is minimizing $-\mathbf{f}$.
- Different from solving $\boldsymbol{\partial f} / \boldsymbol{\partial} \mathrm{x}_{\mathbf{i}}=\mathbf{0}$ because of integration properties.
- Once more, $\mathbf{D}=\mathbf{1}$ is substantially different from higher dimensional cases.
- There's no guarantee that we find the global minimum, methods return local minimums.


## 1D Minimization



- Local minimum: smallest function value in its neighbourhood.
- Global minimum: smallest such function value for all $\mathbf{x}$.


## Strategy

- Generalize bracketing to minimization.
- $\mathbf{a}, \mathbf{b}, \mathbf{c}$ bracket a minimum of $\mathbf{f}$ if

$$
\mathbf{f}(\mathrm{a})>\mathrm{f}(\mathrm{~b})<\mathrm{f}(\mathrm{c})
$$

- Once a minimum is bracketed, we can successively refine the brackets.


## Minimization Methods for 1D

Golden Section Method

- Choose the larger of the intervals [a, b] and [b, c].
- Place a new point d a fraction $\mathbf{w}$ into the larger interval.
- New triplet is succession of points bracketing the minimum.

- Optimal choice for

$$
w=\frac{1}{2}(3-\sqrt{5}) \approx 0.38197
$$

Brent's Method

- Uses a quadratic interpolation using the three points, to find a new point.

All of these are in the GSL library.

## Finding the global minimum

- There's no guarantee that we find the global minimum, methods return local minima.
- If we're not sure: Essentially try again from a different starting points, and check which of the found minima is the lowest.
- Some methods, especially those using random numbers like simulated annealing, have ways of escaping local minima.


## Multidimensional Minimization

- No bracketing method available
- Not as dire as multi-dimensional root finding: down is down, so as local as the function is bounded from below, you're likely to find at least a local minimum.

Common strategy

- Most methods are built upon 1D methods.
- Start from a trial point $\mathbf{P}$
- Perform one dimensional minimizations in D 'independent' directions
- Hope for the best. Test for convergence.


## Finding 'independent’ directions

Just choosing $\mathbf{D}$ directions will lead to staircase scenario.


It might get to the minimum eventually, but requires a lot of small 1d minimizations to do so.

## Finding 'independent' directions

Non-gradient based direction sets

- Use previous 1D minimization to guide the choice of direction sets. Powell's Methods

Gradient based direction sets

- This one uses the derivatives to determine in which D directions to go next.
- Steepest Descent

It may seem like a good idea just to go in the direction of the gradient: steepest descent. But you would have to take many straight turns.

- Conjugate Gradient. Essentially, you're solving the system as if it is quadratic. Better.

All of these are in the GSL library.

## Simulated Annealing

- All of these methods have trouble with functions that have (many) local minima.
- Let's look for inspiration in Nature.
- Why? Nature is a very good minimizer.
- For instance, when temperature drops below $\mathbf{0}^{\circ} \mathrm{C}$, water molecules find a new low energy configuration, which any of our previous minimization methods would have a tough time finding: $\mathbf{D}=\mathcal{O}\left(10^{23}\right)$.
- The mechanism behind this is that while the temperature drops, the molecules are constantly moving in and out of local mimima, thus exploring much more of configurational space.


## Simulated Annealing

Ok, so how does that help us?

- D dimensional space becomes a configuration space.
- f becomes an energy function for the system.
- We assign the system a temperature $\mathbf{T}$.
- We assign the system a dynamics that is such that the equilibrium distribution of configurations is

$$
P(\vec{x}) \propto e^{-f(\vec{x}) / T}
$$

E.g. use the Monte Carlo method from last week's lecture!

- Evolve this system with given dynamics, but from time to time, we reduce the temperature $\mathbf{T}$.
- As $\mathbf{T} \rightarrow \mathbf{0}$ the system will get trapped in a minimum. If we anneal slow enough, there's a good chance it's the global minimum.


## Simulated Annealing


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## Simulated Annealing



## Simulated Annealing



## Simulated Annealing



## Simulated Annealing



## Simulated Annealing

Final notes

- Typically does not converge as fast as the deterministic ones.
- But applicable in many 'hard' cases (travelling salesman).
- GSL has it too.
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