Hybrid OpenMP+MPI

Scientific Computing Lecture 24

SciNet, University of Toronto

Ramses van Zon April 3, 2014



Shared and distributed memory

Modern clusters have a hybrid architecture.

- Multicore machines linked together with an interconnect
- Machines with GPU or other coprocessors: GPU is multi-core, but the amount of shared memory is limited.





MPI vs OpenMP

We have OpenMP for shared memory programming.

We have MPI to program distributed memory machines

model	memory	latency	mem.overhead	scalable	incremental
OpenMP	shared mem	low	low	limited	yes
MPI	distributed	high(er)	higher	yes	no

Could we have the best of both worlds?



Hybrid programming model of using MPI and OpenMP:

- MPI across nodes
- OpenMP within nodes
- Minimizes communication
- Scalable
- Not much more complicated than pure MPI



Pros

- No decomposition on node
- Lower latency, less communication
- Less duplication of data (and perhaps computation)
- OpenMP has load balance capabilities

Cons

- One more layer to maintain
- OpenMP has more hidden side effects
- May have to worry about NUMA (later)



Example

```
#include <mpi.h>
#include <omp.h>
#include <iostream>
int main(int argc, char ** argv)
ł
 int size, rank;
 MPI_Init(&argc,&argv);
 MPI_Comm_get_rank(MPI_COMM_WORLD,&rank);
 MPI_Comm_get_size(MPI_COMM_WORLD,&size);
 #pragma omp parallel for
 for (int i=0;i<4;i++)</pre>
  std::cout << "Hello world from thread "</pre>
              << omp_get_thread_num() << std::endl;
 MPI_Finalize():
}
```





- Memory shared among threads of same process
- Memory not shared among threads of different processes

- Note: OpenMP inside MPI
- Often, one starts with an MPI code and adds in OpenMP.
- Compilation: mpic++ -fopenmp [filename] -o [executable]
- Run: export OMP_NUM_THREADS=M mpirun -np N --bynode [executable]
- This starts N processes, spread out over the nodes.
- Between MPI_Init and MPI_Finalize, each process spawns OMP_NUM_THREADS threads in #pragma omp parallel blocks.



Thread Safety

- Some implementations are limited and do not have support for MPI calls within OpenMP parallel blocks
- Thus, may need to do MPI in serial regions
- Not necessarily bad:





Less communication channels Bigger messages



MPI_Init_thread

... is an MPI_Init replacement that can check for thread support of the MPI implementation.

required and *provided can take values:

MPI_THREAD_SINGLE MPI_THREAD_FUNNELED MPI_THREAD_SERIALIZED MPI_THREAD_MULTIPLE Only 1 thread will execute. 1 thread calls MPI. 1 thread calls MPI at one time. Multiple threads may call MPI at once



Example

```
#include <mpi.h>
#include <omp.h>
#include <iostream>
int main(int argc, char ** argv)
ł
 int size,rank,thread;
 MPI_Init_thread(&argc,&argv,MPI_THREAD_FUNNELED,&thread);
 if (thread>=MPI_THREAD_FUNNELED) {
  MPI_Comm_get_rank(MPI_COMM_WORLD,&rank);
  MPI_Comm_get_size(MPI_COMM_WORLD,&size);
   #pragma omp parallel for
  for (int i=0;i<4;i++)</pre>
    std::cout << "Hello world from thread "</pre>
               << omp_get_thread_num() << std::endl;
 MPI_Finalize();
```



Common useful cases

- Memory bound applications
 - each mpi process is a full application
 - openmp requires less memory
- To fit NUMA (later)
- Overlap comm/comp
 - 1 thread for communication, $i/o, \ services$
 - rest for work







Common useful cases

- Overlap IO/comp
 - 1 thread for IO
 - rest for work







Shared Memory: NUMA

NUMA = Non-Uniform Memory Access

- Multiple cores, but often multiple chips (aka sockets).
- Each chip may have some memory nearby, but can also access the memory of the other chip, at a slower rate.
- Each core typically has some memory/cache of its own.
- Memory locality matters even on a node.





Affinity

Where do processes, threads and memory go?

- Operating system distributes threads and processes over cores, and may migrate them from one core to another.
- Typically, one would want the threads of a process to be on the same core (but not always).
- One would want the memory used by a thread to be close to the core on which it runs.
- ► In Linux, memory is not physically allocated until used.
- Memory is owned by the first thread that uses it: 'first access'.
- Data initialized in a serial section may be 'far' for some threads.

On most systems, these defaults make sense.



Tuning the process/thread affinity

Process/thread affinity:

- Command-line tools like numact1.
- Calls to sched_setaffinity.
- Flags to mpirun.
 E.g. OpenMPI has -bind-to-core and -bind-to-socket.
- OMP_PROC_BIND=true and implementation specific environment variables.

This does prevent the OS from load balancing. Not an issue if you're using all resources of a node.



Tuning the memory affinity

- If a process does not get migrated by OS, memory will remain close to process.
- For threads, in coding, use thread-local variables if you can. (sometimes copying into a thread-local variable can help)
- When a part of the data is mainly used by a specific thread, initialize it in that thread (i.e. not in a serial section).



Homework

- Take your mpi'ed code of last week, and ensure that each processes writes to their own file in binary.
- Change output to MPI-IO so you'd get only one file per run.
- ► Add OpenMP to the loops where it makes sense.
- Set the 'plotEvery' parameter in the diffusion.ini file to 500.
- Time the 16 cases that you get taking values for OMP_NUM_THREADS and the number of MPI processes from the values 1,2,4, and 8.
 Be sure to adjust the number of cores in your submission
 - script to avoid overloading cores.
- Plot the timings as a function of OMP_NUM_THREADS, one plot for each value of the number of MPI processes.
- Compare cases using the same number of overall cores, too.
- Plot the timing results and explain what you see.

Submit code, makefile, git log, plots, explanations by April 1